Secure and Distributed On-Demand Randomized Routing in WSN

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ABSTRACT

Security and energy efficiency is of paramount importance in a wireless sensor network. This is due to their vulnerable deployment conditions and battery based power. This paper presents a secure and distributed algorithm that generates routes on-demand in a wireless sensor network. Dynamic route generation is facilitated by PSO, a metaheuristic technique. Current network traffic in that route and charge contained in the candidate node are used as evaluation parameters along with the node distance, hence a huge reduction in the packet loss was observed. Experiments were conducted and it was observed that the proposed algorithm exhibits very low selection overhead and also provides distributed routs, which eventually lead to prolonged network lifetime.

Indexing terms/Keywords

Altruism; Particle Swarm Optimization; Routing; Security; Selfishness; Wireless Sensor Networks

INTRODUCTION

Wireless Sensor Networks (WSN) have been very widely used in the recent times and has attracted several researches into this domain [13]. With the advancement in technology, WSNs are being used in several areas such as military [14], environmental applications [15], locating forest fires [16], observing animal habitats [17] and health care [18]. WSNs, though being very useful, are prone to failures and are vulnerable. The most important vulnerability arises from the fact that they are deployed mostly in inaccessible areas, hence cannot be frequently monitored. Once deployed, they function autonomously. Hence, they only transfer their own packets and forwarding packets of other nodes. Every transmission causes a reduction in charge in the node. As the network operates, some nodes are operated on more than others, hence charge in certain nodes begins to deplete faster. This leads to the nodes turning selfish. Selfish nodes do not participate in packet forwarding, instead they only transfer their own packets. This behavior is incorporated in-order to retain the node from completely getting depleted.

The above mentioned properties are specific to WSNs and hence a routing algorithm designed for WSN should consider these properties and not just the shortest routing path to determine routes. There always exists a tradeoff between security and energy efficiency. This tradeoff is determined by the type of application the WSN is being designed to be operated on. This paper presents a randomized routing technique using PSO that generates dynamic random routes during transmissions. The routes are usually distributed throughout the network, hence specific node based charge depletions are avoided. Further, the dynamic on-demand routes provides security to the transmitted data.

The remainder of this paper is structured as follows; section II provides the related works, section III presents an elaborate discussion on the proposed approach, section IV presents the results and section V concludes the study.

RELATED WORKS

Wireless Sensor Networks (WSN), being a widely used technology has several contributions dealing with providing improvements in transmission security. Some of the recent contributions are discussed below.

A Bayesian Signaling Game model is presented in [1] that analyzes the strategy profiles for normal and anomalous nodes. This method also involves payoffs to motivate the nodes in abstaining them from malicious behavior. This is a variant of the trust based approach, where every node identifies and maintains the belief levels of their neighbors. These levels are periodically updated to provide effective results. The downsides of this approach is that it has high computational complexities, which makes them suitable only for the high performance based networks. Similar profiling based methods include [8]. Contributions by Roy et al. [2] served as the base for game theory based approaches. This method analyzed the usage of game theory on networked applications. Approaches that can be used to apply game theory concepts in network security was presented in [3].
A scalable and secure routing technique that uses the received signal strength in WSNs to perform effective routing is presented in [4]. A distributed location verification algorithm is used to identify the signal strength. This helps avoid spoofing attacks to a large extent. Scalability is provided by utilizing the broadcast nature of the WSN. The ambient trust model is also incorporated in this architecture to avert several other attacks. Several other location based techniques [5-7] are available that avert spoofing attacks using the signal strength of the received packets to identify the location of the transmitting node.

A node habit based routing technique is presented in [12]. This method analyzes and records the behavior of nodes in a network. These behavior patterns are used to perform the node selection process. Privacy of the nodes are preserved using cryptographic techniques. This is an on-demand routing algorithm that best operates on delay tolerant data. other on-demand routing techniques are discussed in [9-11].

SECURE AND DISTRIBUTED ON-DEMAND RANDOMIZED ROUTING IN WSN

Security and energy efficiency is of paramount importance in a wireless sensor network. The proposed approach presents an on-demand routing technique that uses randomized routing to provide security. Further, due to the process of randomized node selection, this method provides effective distribution of load, hence increasing the lifetime of the network. Figure 1 presents the architecture of the proposed methodology.

A wireless sensor network is a dynamic structure that must be flexible in order to accommodate network structure changes. Sensor nodes are prone to failures, hence the network details must be frequently updated. Since the proposed approach for routing requires the details of neighbor nodes, a hello packet is transmitted to all the 1-hop neighbors and their acknowledgements are collected to update the current live node list in the network and their current position. It becomes mandatory for every node to maintain the updated list in-order to avoid failed transmissions, since failed transmissions tend to be much costlier in WSNs when compared to normal networks.

On initiation of transmission, the packets are constructed and PSO based routing is initialized by the architecture. PSO [19], being a metaheuristic algorithm is a valid candidate for providing distributed routes. The randomness involved in PSO provides effective random routes that aids in the uniform use of energy in the network rather than usage of certain specific nodes. Another major problem in WSNs tend to be lack of memory and computational resources. PSO, when compared to other meta-heuristic techniques tend to utilize lesser memory and also has a low computational overhead [20].

On obtaining a transmission initiation signal, PSO builds the search space. The search space of PSO is composed of the current node and all the neighbor nodes and their corresponding properties [21]. This method considers the current network traffic recorded in that route and the current charge of the node. Utilizing a route with heavy traffic will often lead to transmission failures, while opting for a node with low charge will lead to the node turning selfish, which has a bad effect on the network’s overall functioning. These two properties remain to be the most important properties governing safe transmission and in maintaining network stability. Hence they are used as components during the node selection phase.

All the required particles are distributed on the current node and the particle best (pbest) and global best (gbest) values are set to the current node. The initial velocity of a node must be assigned in-order to initiate movement in the search space. The initial velocity calculation is performed in a random manner and is maintained to be within the bounds of the current search space.

\[ V_{i,d} \sim U(-|b_{up} - b_{lo}|, |b_{up} - b_{lo}|) \]  

where \( b_{lo} \) and \( b_{up} \) are the lower and upper bounds of the search space respectively.

Particle acceleration is then triggered and the particles start their movement using the initial velocity and direction obtained from the previous phase. After the initial movement, the particles are distributed in the search space. PSO is continuous in nature, while the current application demands a discrete node as a solution. Hence a discretization function is used to identify the final node.

\[ P' = \min \left( \sum_{j=1}^{p} \left( \sum_{k=1}^{N} \sqrt{(P_{ik} - N_{k})^2} \right) \right) \text{ for } i = 1 \text{ to } p \]  

where \( P_0 \) refers to the particle \( i \)‘s current location corresponding to dimension \( k \), \( N_k \) refers to the \( k \)th dimension of node \( N \).

This process is repeated for all particles and further movements are triggered by

\[ V_{i,d} \leftarrow \omega V_{i,d} + \varphi_p P_{i,d} - X_{i,d} \] \[ + \varphi_g P_{g,d} - X_{g,d} \]  

where \( \omega \), \( \varphi_p \) and \( \varphi_g \) are the random numbers, \( P_{i,d} \) and \( g_{d} \) are the parameter best and the global best values, \( X_{i,d} \) is the value current particle position, and the parameters \( \omega, \varphi_p, \) and \( \varphi_g \) are selected by the practitioner.

The gbest obtained after satisfying the termination criterion is considered as the next probable node for transmission of the packet. This process is repeated until the packet reaches the destination node.
Fig 1. Secure and Distributed On-Demand Randomized Routing in WSN – Architecture
RESULTS AND DISCUSSION

Experiments were conducted on a network and the efficiency of the algorithm was measured in terms of selection overhead and load distribution among the nodes.

![Selection Overhead](image1)

**Fig 2. Selection Overhead**

Fig 2 shows the selection overhead incurred in the path selection process. It shows the time taken to identify a complete path in the network for traversing all the nodes. It could be observed that the average time taken to traverse all the nodes in the network is approximately 6.5ms.

![Distance Covered](image2)

**Fig 3. Total Path Covered**

Fig 3 shows the total distance covered by the algorithm while traversing all the nodes in the network. It was observed that a minimum distance of 470 and a maximum distance of 565 was taken by the algorithm. This increase in the distance is attributed to the load distribution component in the algorithm.

![Selection Overhead](image3)

**Fig 4. Selection Overhead for Specific Path**

Fig 4 presents the selection overhead for the algorithm to traverse between two specific nodes. It was observed that a maximum overhead of 6ms and a minimum overhead of 0ms was observed.
Fig 5. Distance Covered for Specific Path

Fig 5 shows the distance covered by the algorithm for traversing a specific path. It was observed that the distance covered ranges from 50 to 370. This is also attributed to the dispersion mechanism in the algorithm to help in load distribution and avoid selfish nodes in the network.

Fig 6. Node Usage Levels (Load Distribution)

Fig 6 shows the load distribution in the network for obtaining a path between specific nodes. The highest path access has occurred in the starting and ending nodes, while other nodes exhibit an average access throughout. For this experiment, the start and the end nodes are set to node 2 and node 4 respectively. Hence high access is observed in those nodes, while other nodes exhibit an average access. Hence it could be concluded that the proposed algorithm functions well in the distribution of load.

CONCLUSION

Several techniques exist in literature to perform secure communications in WSN. Some of these techniques also incorporate energy efficiency, but most of them are application specific. This paper presents a generic approach that provides both the major requirements of WSNs. The on-demand and dynamic nature of the algorithm has proved to be a major advantage when it comes to security and in maintaining the network’s stability. Future extensions of this work includes utilizing a hybrid technique to reduce the selection overhead further. Though the selection overhead remains low, reducing it further will improve the ability of the algorithm to support real time traffic in the network. Further, trust levels of nodes can be recorded and incorporated into the algorithm to aid in the process of node selection, which can help prevent transmission failures to a large extent.
REFERENCES


Author’s biography with Photo

Mr. V. Upendran is pursuing his Ph.D in computer Science from Bharathiar University, Tamil Nadu, India. He is currently Head, Department of Computer Science, Srimad Andavan Arts and Science College (Autonomous), Affiliated to Bharathidasan University, Trichirappalli, India. He has 10 years of teaching experience. He has presented 2 International and 4 National papers in conferences. He has published 4 papers in International and one in National journal.

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